

数理情報学専攻

修士課程入学試験問題

専門科目 数理情報学

平成21年8月25日(火) 10:00 ~ 13:00

5問出題, 3問解答

- This booklet is an informal English translation of the original examination booklet.

- **Answer three problems out of Problem 1 ~ Problem 5.**
- **Answer in Japanese or English.**

Problem 1

- (1) Show that for a positive-definite symmetric real matrix A there uniquely exists a positive-definite symmetric real matrix B such that $A = B^2$.
- (2) Show that for a nonsingular real matrix F there uniquely exist an orthogonal matrix R and a positive-definite symmetric real matrix U such that $F = RU$.
- (3) Show that for a nonsingular real matrix F there uniquely exist an orthogonal matrix R and positive-definite symmetric real matrices U and V such that $F = RU = VR$.

Problem 2

Each item produced in a factory is defective with probability ε ($0 < \varepsilon < 1$). Let X be a random variable such that $X = 0$ if the item is non-defective and $X = 1$ if it is defective.

Answer the following questions (1), (2) and (3).

- (1) Each produced item is inspected before shipment. Let Y be a random variable such that $Y = 0$ if the inspection result indicates that the item is non-defective and $Y = 1$ if the inspection result indicates that the item is defective. When $X = 0$, we have $Y = 0$ with probability 0.9 and $Y = 1$ with probability 0.1. When $X = 1$, we have $Y = 0$ with probability 0.1 and $Y = 1$ with probability 0.9. The item is shipped if $Y = 0$, and it is discarded if $Y = 1$. Obtain the probability $P(X = 1|Y = 0)$ that an item shipped from the factory is defective.
- (2) The same inspection as in (1) is repeated n times before the produced item is shipped. Let Z_n be a random variable to represent the number of inspection results indicating that the item is defective. Here the n inspection results are independent of each other. Show that the conditional probability $P(X = 1|Z_n = z_n)$ ($z_n = 0, 1, \dots, n$) is a function of ε and $n - 2z_n$, the latter of which is equal to the difference between the number $n - z_n$ of inspection results indicating that the item is non-defective and the number z_n of inspection results indicating that the item is defective.
- (3) Let $q = P(X = 1|Z_2 = 0)$ and $r = P(X = 1|Z_2 = 2)$. The same inspection as in (1) is repeated until the conditional probability $P(X = 1|Z_n = z_n)$ based on the inspection result $Z_n = z_n$ first becomes at most q or at least r . Obtain the expected number of inspections performed before termination when the item is non-defective.

Problem 3

- (1) Suppose that A is a bounded set[†] in an n -dimensional Euclidean space and that the volume $v(A)$ of A is greater than 1 (i.e., $v(A) > 1$). We consider the intersection $(A + \mathbf{g}) \cap C$ of

$$A + \mathbf{g} = \{\mathbf{x} + \mathbf{g} \mid \mathbf{x} \in A\},$$

which denotes the set obtained by translating A by an n -dimensional integral vector \mathbf{g} , and

$$C = \{\mathbf{x} \mid \mathbf{x} = (x_1, \dots, x_n), 0 \leq x_i < 1 \ (i = 1, \dots, n)\}.$$

Prove

$$\sum_{\mathbf{g} \in \mathbb{Z}^n} v((A + \mathbf{g}) \cap C) = v(A),$$

where \mathbb{Z}^n denotes the set of all n -dimensional integral vectors.

- (2) Prove that there exist two distinct integral vectors \mathbf{g} and \mathbf{h} such that the intersection of $(A + \mathbf{g}) \cap C$ and $(A + \mathbf{h}) \cap C$ is nonempty.
- (3) Prove that there exist two distinct points \mathbf{x} and \mathbf{y} in A such that $\mathbf{x} - \mathbf{y} \in \mathbb{Z}^n$.
- (4) Suppose that B is a bounded convex set[‡] which is symmetric with respect to the origin in the n -dimensional Euclidean space and satisfies $v(B) > 2^n$. Prove that B contains a nonzero integral vector by applying the result of (3) to the set $\frac{1}{2}B = \{\frac{1}{2}\mathbf{x} \mid \mathbf{x} \in B\}$.
- (5) Let $R = (r_{ij})$ be a nonsingular real matrix of order 3, and denote the determinant of R by $\det R$. Prove that for any α satisfying the inequality $\alpha > \sqrt[3]{3!|\det R|}$ there exists a nonzero integral vector $\mathbf{g} = (g_1, g_2, g_3) \in \mathbb{Z}^3$ such that

$$\sum_{i=1}^3 \left| \sum_{j=1}^3 r_{ij} g_j \right| < \alpha.$$

[†] Strictly speaking, A is a bounded measurable set, and $v(A)$ is the measure of A .

[‡] B is said to be symmetric with respect to the origin, if $\mathbf{x} \in B$ implies $-\mathbf{x} \in B$. B is said to be convex if, for any two points \mathbf{x} and \mathbf{y} in B , it contains the line segment joining \mathbf{x} and \mathbf{y} . In particular, for any two points \mathbf{x} and \mathbf{y} in B , we have $\frac{1}{2}(\mathbf{x} + \mathbf{y}) \in B$. It is known that a bounded convex set is measurable.

Problem 4

Let us consider a system consisting of three water-tanks 1, 2, and 3 (Fig. 1).

Water in the three tanks constantly flows out at the rates of r_1 , r_2 , and r_3 per unit time, respectively. At the same time, water is supplied to one of the three tanks at the rate of $R = r_1 + r_2 + r_3$ per unit time. When a tank becomes empty, the water supply is switched to the tank. Assume that the system stops if two tanks simultaneously become empty. Water does not overflow from the tanks.

Let $q(t) \in \{1, 2, 3\}$ denote the tank to which water is supplied at time t . Then, the amount of water $w_p(t)$ in tank p at time t obeys the following differential equation:

$$\frac{dw_p(t)}{dt} = R\delta_{pq(t)} - r_p \quad (p = 1, 2, 3),$$

where δ is the Kronecker delta ($\delta_{pq} = 1$ if $p = q$ and $\delta_{pq} = 0$ if $p \neq q$).

In the following, assume that $r_1 = r_2 = r_3 = \frac{1}{3}$ and that the total amount of water in the three tanks $w_1(t) + w_2(t) + w_3(t)$, which is constant, is equal to 1.

- (1) Let t_i denote the time when one or two tanks become empty for the i th time after time 0. If two tanks simultaneously become empty and the system stops at time t_i , the subsequent sequence is defined by $t_i = t_{i+1} = t_{i+2} = \dots$.

Suppose that tank 1 becomes empty at time t_i . Let $\alpha = w_3(t_i)$ be the amount of water in tank 3 at time t_i . Express the amounts of water in the three tanks $w_1(t_{i+1})$, $w_2(t_{i+1})$, and $w_3(t_{i+1})$ at time t_{i+1} using α .

- (2) Since one or two tanks are empty at time t_i , the state of the system at that moment can be described by a point x_i in the half-open interval $[0, 1) = \{x \mid 0 \leq x < 1\}$ as follows:

$$x_i = \begin{cases} \frac{w_3(t_i)}{3} & \text{(if } w_1(t_i) = 0 \text{ and } w_2(t_i) \neq 0), \\ \frac{w_1(t_i) + 1}{3} & \text{(if } w_2(t_i) = 0 \text{ and } w_3(t_i) \neq 0), \\ \frac{w_2(t_i) + 2}{3} & \text{(if } w_3(t_i) = 0 \text{ and } w_1(t_i) \neq 0). \end{cases}$$

Then, x_{i+1} is uniquely determined from x_i , and can be written in the form of $x_{i+1} = f(x_i)$. Obtain the map $f: [0, 1) \rightarrow [0, 1)$. Draw the graph of f .

- (3) Show that for any two distinct points x and $y \in [0, 1)$, there exists an integer $k \geq 0$ such that $|f^k(y) - f^k(x)| \geq \frac{1}{4}$. Here $f^0(x) = x$ and $f^k(x) = f(f^{k-1}(x))$ ($k = 1, 2, \dots$).

[Hint] First show that, if $|y - x| < \frac{1}{4}$, then $|f(y) - f(x)| \geq 2|y - x|$.

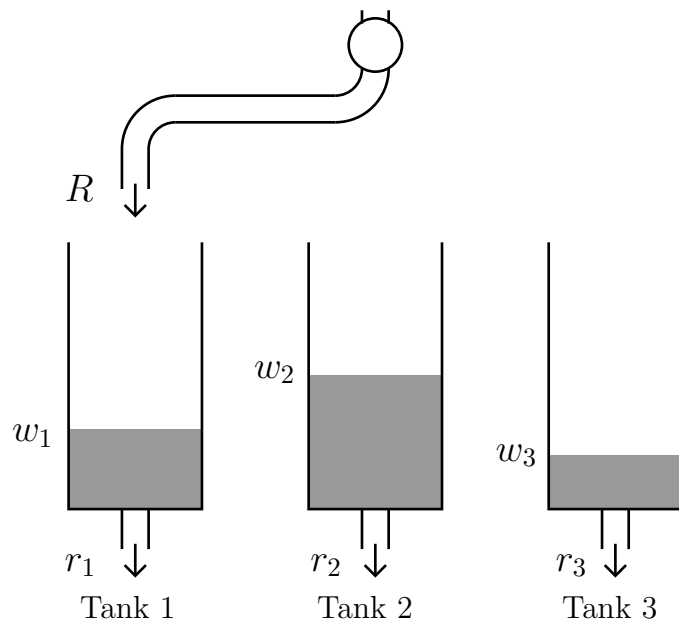


Fig. 1

Problem 5

- (1) Given $2n+1$ positive integers $p_1, p_2, \dots, p_n, s_1, s_2, \dots, s_n, S$, let us consider solving the following optimization problem by the dynamic programming method.

$$\begin{aligned}
 \text{(Problem P)} \quad & \max \sum_{i=1}^n p_i x_i \\
 & \text{subject to } \sum_{i=1}^n s_i x_i \leq S \quad (*) \\
 & x_i \in \{0, 1\} \quad (i = 1, \dots, n). \quad (**)
 \end{aligned}$$

In order to solve Problem P, for $j = 1, \dots, n$ and $s = 0, 1, \dots, S$, we consider the optimization problem

$$\begin{aligned}
 \max \quad & \sum_{i=1}^j p_i x_i \\
 \text{subject to} \quad & \sum_{i=1}^j s_i x_i = s \\
 & x_i \in \{0, 1\} \quad (i = 1, \dots, j),
 \end{aligned}$$

and denote by $A(j, s)$ the optimal value (i.e., the maximum value of the objective function) of this problem. Here we define $A(j, s) = -\infty$ if the problem has no feasible solution. Then we can solve Problem P by letting $A(0, 0) = 0$ and $A(0, s) = -\infty$ ($s \neq 0$), and by using the following recurrence for $A(j, s)$ ($j = 1, \dots, n$):

$$\begin{aligned}
 A(j, s) &= A(j-1, s), \quad \text{if } s < s_j, \\
 A(j, s) &= \max\{A(j-1, s), p_j + A(j-1, s - s_j)\}, \quad \text{otherwise.}
 \end{aligned}$$

Compute the optimal value for the problem instance with $n = 5$, $S = 5$, $p_1 = 2$, $p_2 = 3$, $p_3 = 2$, $p_4 = 1$, $p_5 = 3$, $s_1 = 2$, $s_2 = 3$, $s_3 = 1$, $s_4 = 2$, $s_5 = 1$ by applying the dynamic programming method based on the recurrence above. Give all the values $A(j, s)$ which are needed to compute the optimal value.

- (2) Give an algorithm based on the recurrence in (1). Show the time complexity of the algorithm and explain whether it is polynomial in the input length.
- (3) Consider Problem Q that is obtained from Problem P by replacing constraint (**) with $x_i \in \{0, 1, \dots, 10\}$ ($i = 1, \dots, n$). Obtain a recurrence as in (1).
- (4) Consider Problem R that is obtained from Problem P by adding a new constraint $\sum_{i=1}^n w_i x_i \leq W$, where w_1, \dots, w_n , and W are all positive integers. Obtain a recurrence as in (1).